



Multiagent net and Algebraic connectivity

Ya-Lei Jin (晋亚磊)

Shanghai Jiao Tong University

yaleijin@sjtu.edu.cn

16th Nov. 2013

This work is joined with my advisor Xiao-dong Zhang (张晓东)





Outline

- Introduction .
- The multiagent net.
- The algebraic connectivity.
- The main theorem.



Outline

- Introduction .
- The multiagent net.
- The algebraic connectivity.
- The main theorem.



Outline

- Introduction .
- The multiagent net.
- The algebraic connectivity.
- The main theorem.



Outline

- Introduction .
- The multiagent net.
- The algebraic connectivity.
- The main theorem.



Definition and Background

- Let $G = (V, E)$ be a simple graph (without loop and multiedges) with vertex set $V = [n]$ and edge set $E(G)$.
- $A(G) = (a_{ij})$: the adjacency matrices of G with $a_{ij} = 1$ if $v_i \sim v_j$; $a_{ij} = 0$ otherwise.
- $D(G) = \text{diag}(d_1, \dots, d_n)$ degree diagonal matrix with d_i being the degree of vertex v_i .



Definition and Background

- Let $G = (V, E)$ be a simple graph (without loop and multiedges) with vertex set $V = [n]$ and edge set $E(G)$.
- $A(G) = (a_{ij})$: the adjacency matrices of G with $a_{ij} = 1$ if $v_i \sim v_j$; $a_{ij} = 0$ otherwise.
- $D(G) = \text{diag}(d_1, \dots, d_n)$ degree diagonal matrix with d_i being the degree of vertex v_i .



Definition and Background

- Let $G = (V, E)$ be a simple graph (without loop and multiedges) with vertex set $V = [n]$ and edge set $E(G)$.
- $A(G) = (a_{ij})$: the adjacency matrices of G with $a_{ij} = 1$ if $v_i \sim v_j$; $a_{ij} = 0$ otherwise.
- $D(G) = \text{diag}(d_1, \dots, d_n)$ degree diagonal matrix with d_i being the degree of vertex v_i .



- The *Laplacian matrix* of G : $L(G) = D(G) - A(G)$.
- $L(G)$ is positive semi-definite and singular.
- The eigenvalues of $L(G)$ are denoted by

$$\lambda_1(G) \geq \lambda_2(G) \geq \cdots \geq \lambda_{n-1}(G) \geq \lambda_n(G) = 0;$$

called the *Laplacian eigenvalues* of G . In particular, $\lambda_{n-1}(G)$ is called the algebraic connectivity of G and denoted by $\alpha(G)$.

- $\mathbf{1}$ is an eigenvector corresponding $\lambda_n(G)$.



- The *Laplacian matrix* of G : $L(G) = D(G) - A(G)$.
- $L(G)$ is positive semi-definite and singular.
- The eigenvalues of $L(G)$ are denoted by

$$\lambda_1(G) \geq \lambda_2(G) \geq \cdots \geq \lambda_{n-1}(G) \geq \lambda_n(G) = 0;$$

called the *Laplacian eigenvalues* of G . In particular, $\lambda_{n-1}(G)$ is called the algebraic connectivity of G and denoted by $\alpha(G)$.

- $\mathbf{1}$ is an eigenvector corresponding $\lambda_n(G)$.



- The *Laplacian matrix* of G : $L(G) = D(G) - A(G)$.
- $L(G)$ is positive semi-definite and singular.
- The eigenvalues of $L(G)$ are denoted by

$$\lambda_1(G) \geq \lambda_2(G) \geq \cdots \geq \lambda_{n-1}(G) \geq \lambda_n(G) = 0;$$

called the *Laplacian eigenvalues* of G . In particular, $\lambda_{n-1}(G)$ is called the algebraic connectivity of G and denoted by $\alpha(G)$.

- $\mathbf{1}$ is an eigenvector corresponding $\lambda_n(G)$.



- The *Laplacian matrix* of G : $L(G) = D(G) - A(G)$.
- $L(G)$ is positive semi-definite and singular.
- The eigenvalues of $L(G)$ are denoted by

$$\lambda_1(G) \geq \lambda_2(G) \geq \cdots \geq \lambda_{n-1}(G) \geq \lambda_n(G) = 0;$$

called the *Laplacian eigenvalues* of G . In particular, $\lambda_{n-1}(G)$ is called the algebraic connectivity of G and denoted by $\alpha(G)$.

- $\mathbf{1}$ is an eigenvector corresponding $\lambda_n(G)$.



Agreement protocol(**):

For a graph $G = (V, E)$, $V = [n]$, we define a agreement protocol as the following:

$$\dot{x}(t) = -L(G)x(t),$$

i.e.

$$\dot{x}_i(t) = \sum_{j \in N(i)} (x_j(t) - x_i(t)),$$

where $x_i(t)$ is a value at the time t and at the vertex i .



Three Problems:

- 1. Whether $x(t)$ is convergent?
- 2. If $x(t)$ is convergent, then what is the form of $x(t)$ if t tends infinity.
- 3. If $x(t)$ is convergent, then what is the convergent rate?



Three Problems:

- 1. Whether $x(t)$ is convergent?
- 2. If $x(t)$ is convergent, then what is the form of $x(t)$ if t tends infinity.
- 3. If $x(t)$ is convergent, then what is the convergent rate?



Three Problems:

- 1. Whether $x(t)$ is convergent?
- 2. If $x(t)$ is convergent, then what is the form of $x(t)$ if t tends infinity.
- 3. If $x(t)$ is convergent, then what is the convergent rate?



- Using the spectral factorization of the Laplacian, one has

$$e^{-L(G)t} = e^{-\lambda_1(G)t} u_1 u_1^T + \dots + e^{-\lambda_{n-1}(G)t} u_{n-1} u_{n-1}^T + e^{-\lambda_n(G)t} u_n u_n^T.$$

where u_i is an eigenvector corresponding to λ_i and $u_i u_j^T = 0$,
 $j \neq i = 1, 2, \dots, n$.

- The solution of (**), initialized from $x(0) = x_0$, is
 $x(t) = e^{-L(G)t} x_0$.
- Let $\mathcal{A} = \{x \in R^n | x_i = x_j, \text{ for all } i, j\}$



- Using the spectral factorization of the Laplacian, one has

$$e^{-L(G)t} = e^{-\lambda_1(G)t} u_1 u_1^T + \dots + e^{-\lambda_{n-1}(G)t} u_{n-1} u_{n-1}^T + e^{-\lambda_n(G)t} u_n u_n^T.$$

where u_i is an eigenvector corresponding to λ_i and $u_i u_j^T = 0$,
 $j \neq i = 1, 2, \dots, n$.

- The solution of (**), initialized from $x(0) = x_0$, is
 $x(t) = e^{-L(G)t} x_0$.
- Let $\mathcal{A} = \{x \in R^n | x_i = x_j, \text{ for all } i, j\}$



- Using the spectral factorization of the Laplacian, one has

$$e^{-L(G)t} = e^{-\lambda_1(G)t} u_1 u_1^T + \dots + e^{-\lambda_{n-1}(G)t} u_{n-1} u_{n-1}^T + e^{-\lambda_n(G)t} u_n u_n^T.$$

where u_i is an eigenvector corresponding to λ_i and $u_i u_j^T = 0$,
 $j \neq i = 1, 2, \dots, n$.

- The solution of (**), initialized from $x(0) = x_0$, is
 $x(t) = e^{-L(G)t} x_0$.
- Let $\mathcal{A} = \{x \in R^n | x_i = x_j, \text{ for all } i, j\}$



Theorem 1

*Let G be a connected graph. Then the agreement protocol (**)
converges to the agreement set \mathcal{A} with a rate of convergence that
is dictated by $\lambda_{n-1}(G)$.*



Spectral Turán Theorem

Theorem 2 (Jin-Zhang)

Let G be a non-complete graph of order n not containing K_{r+1} .

Then

$$\alpha(G) \leq n - \lceil \frac{n}{r} \rceil = \alpha(T_{n,r}), \quad (1)$$

where $\lceil a \rceil$ is the least integer no less than a . Moreover, if $n = kr$ or $n = kr + r - 1$, then equality (1) holds if and only if G is Turán graph $T_{n,r}$. If $n = kr + t, 0 < t < r - 1$, then equality (1) holds if and only if there exist graphs H_1, \dots, H_t of order $k + 1$ with no edges and H of order $n - (k + 1)t$ not containing K_{r+1-t} such that

$$G = H_1 \vee H_2 \cdots \vee H_t \vee H$$

and $\alpha(H) \geq n - (k + 1)(t + 1)$.



In the following, we give the graphs given clique number with minimum algebraic connectivity:

Theorem 3

Let G be a connected graph with the clique number $r \geq 2$. Then

$$\alpha(G) \geq \alpha(Ki_{n,r}), \quad (2)$$

where $Ki_{n,r}$ is a kite graph of order n which is obtained by adding a pendant path of length $n - r$ to a vertex of K_r . Moreover, equality (2) holds if and only if $G = Ki_{n,r}$.



Proof of the Theorem 3

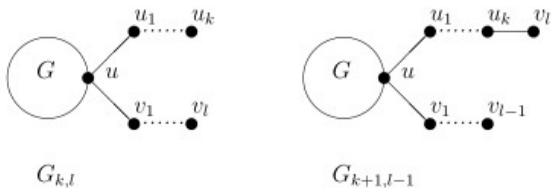


Fig. 1



Proof of the Theorem 3

Lemma 4 (Guo 2010)

Let G be a connected graph with at least two vertices and two paths $P : uu_1u_2 \dots u_k$ and $Q : uv_1 \dots v_l$ of lengths k, l ($k \geq l \geq 1$). If $G_{k,l}$ is the graph obtained from G by attached two paths P, Q at vertex u and $G_{k+1,l-1}^{(1)} = G_{k,l} - v_{l-1}v_l + u_kv_l$ (see Fig. 1) and $k \geq l \geq 1$, then

$$\alpha(G_{k,l}) \geq \alpha(G_{k+1,l-1}^{(1)}). \quad (3)$$

Moreover, inequality (3) is strict if either $X(v_1) \neq 0$ or $X(u_1) \neq 0$.



Proof of the Theorem 3

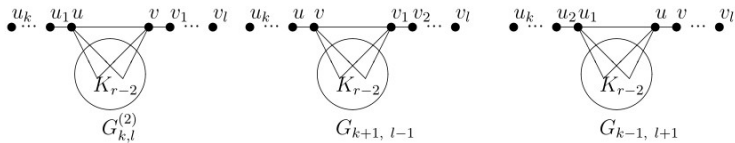


Fig.3



Proof of the Theorem 3

Lemma 5 (Jin-Zhang)

Let $G_{k,l}^{(2)}$ be the graph obtained from a complete graph K_r , $r \geq 3$ with vertex set $V(K_r) = \{w_1, \dots, w_{r-2}, u, v\}$ by attached two paths $P = uu_1 \dots u_k$ and $Q = vv_1 \dots v_l$ at vertices u and v , respectively (see Fig. 3). If

$$G_{k+1,l-1}^{(2)} = G - \{uw_i, 1 \leq i \leq r-2\} + \{v_1w_i, 1 \leq i \leq r-2\},$$

$$G_{k-1,l+1}^{(2)} = G - \{vw_i, 1 \leq i \leq r-2\} + \{u_1w_i, 1 \leq i \leq r-2\}.$$

Then

$$\alpha(G_{k,l}^{(2)}) > \min\{ \alpha(G_{k+1,l-1}^{(2)}), \alpha(G_{k-1,l+1}^{(2)}) \}. \quad (4)$$



Proof of the Theorem 3

Corollary 6 (Jin-Zhang)

Let $G_{k,l}^{(2)}$ be the graph obtained from a complete graph K_r , $r \geq 3$ with vertex set $V(K_r) = \{w_1, \dots, w_{r-2}, u, v\}$ by attached two paths $P = uu_1 \dots u_k$ and $Q = vv_1 \dots v_l$ at vertices u and v , respectively. If $k \geq l \geq 1$, then

$$\alpha(G_{k,l}^{(2)}) > \alpha(G_{k+1,l-1}^{(2)}). \quad (5)$$

Proof of the above Corollary:

If $k = l$, the assertion follows from Lemma 5 and

$G_{k+1,l-1}^{(2)} = G_{k-1,l+1}^{(2)}$. If $k > l$, suppose that $\alpha(G_{k,l}^{(2)}) > \alpha(G_{k+1,l-1}^{(2)})$ does not hold, i.e., $\alpha(G_{k,l}^{(2)}) \leq \alpha(G_{k+1,l-1}^{(2)})$. Then by Lemma 5, we have

$$\alpha(G_{k,l}^{(2)}) > \min\{\alpha(G_{k-1,l+1}^{(2)}), \alpha(G_{k+1,l-1}^{(2)})\} = \alpha(G_{k-1,l+1}^{(2)}). \quad (6)$$

By repeated uses of Lemma 5 and (6), we have

$$\alpha(G_{k-1,l+1}^{(2)}) > \min\{\alpha(G_{k-2,l+2}^{(2)}), \alpha(G_{k,l}^{(2)})\} = \alpha(G_{k-2,l+2}^{(2)}). \quad (7)$$

By repeated uses of Lemma 5, we obtain

$$\alpha(G_{k,l}^{(2)}) > \alpha(G_{k-1,l+1}^{(2)}) > \alpha(G_{k-2,l+2}^{(2)}) > \cdots > \alpha(G_{k-(k-l),l+(k-l)}^{(2)}) = \alpha(G_{l,k}^{(2)})$$

which is a contradiction, since $G_{k,l}^{(2)} = G_{l,k}^{(2)}$. So the assertion holds.



Proof of the Theorem 3

Corollary 7 (Jin-Zhang)

Let $G_{k,l}^{(2)}$ be a graph of order n obtained from the complete graph K_r , $r \geq 3$ with vertex set $V(K_r) = \{w_1, \dots, w_{r-2}, u, v\}$ by attached two paths $P = uu_1 \dots u_k$ and $Q = vv_1 \dots v_l$ at vertices u and v , respectively, where $n = r + k + l$. If $k > 0, l > 0$, then

$$\alpha(G_{k,l}^{(2)}) > \alpha(Ki_{n,r}). \quad (8)$$



Corollary 8

Let G be a graph of order n with clique number r . Then

$$\frac{n}{n - \alpha(G)} \leq r \leq n + 1 - \frac{4}{n\alpha(G)}.$$



上海交通大学

Shanghai Jiao Tong University

Thanks For your Attentions!

